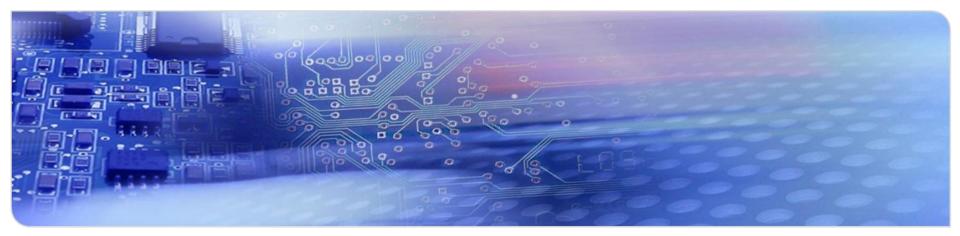




Communication Systems and Protocols

Exercise: FireWire



IEEE Std 1394-2008

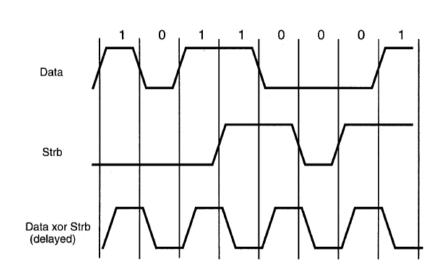


- Three primary applications have driven the design and architecture of the serial bus:
 - an alternate for a parallel backplane bus
 - a low-cost peripheral bus and
 - a bus bridge between architecturally compatible 32-bit buses.
- Version:
 - The description provided here is based on IEEE Std 1394-2008

Data Transmission: Data-Strobe (DS) encoding



- During packet transmission, there is only a single node transmitting on the bus, so the entire media can operate in a half-duplex mode using two signals: Data and Strb.
- NRZ data are transmitted on Data and is accompanied by the Strb signal, which changes state whenever two consecutive NRZ data bits are the same, ensuring that a transition occurs on either Data or Strb for each data bit.
- A clock that transitions each bit period can be derived from the exclusive-or of Data with Strb as shown in Figure

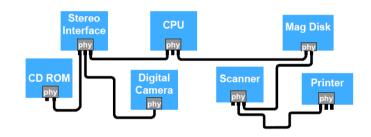


DS Encoding

Cable environment



- Here we focus on the arbitration in the cable environment.
- First describe the cable environment in detail:
 - The cable environment is a network of nodes connected by point-to-point links called physical connections.
 - The physical connection consists of a port on the PHY of each node and the cable between them.
 - The primary restriction is that nodes have to be connected together as an acyclic graph (no loops).



Cable Environment

Fig Source: https://grouper.ieee.org/groups/802/802_tutorials/04-July/NewTechIntroTo1394.pdf

System Configuration

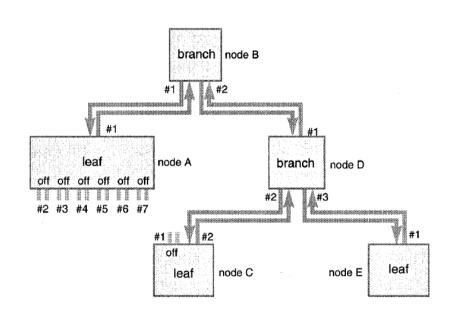


- The cable arbitration takes advantage of the point-to-point nature of the cable environment by having each node handshake with its immediate neighbors to determine ownership of the media.
- Prior to normal operation the system has to be configured.
- System configuration is done in three phases
 - 1. Bus initialize
 - 2. Tree identify
 - 3. Self-identify.

1. Bus Initialize



- Whenever a node joins the bus, a bus reset signal forces all nodes into a special state that clears all topology information.
- After the bus initialization process, the only information known to a node is whether it is a
 - Branch: more than one directly connected neighbor or
 - Leaf: only a single neighbor or
 - Isolated: unconnected
- The eventual root may be either a branch or a leaf!



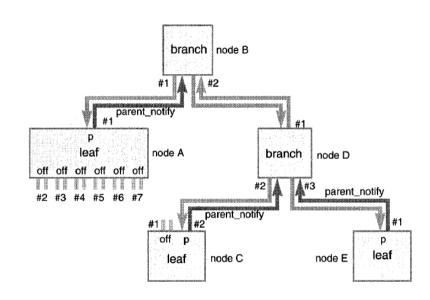
After Bus Initialization



- The tree identify process translates the general network topology into a tree, where one node is designated a root and all of the physical connections have a direction associated with them pointing towards the root node.
- The direction is set by labeling each connected port as a
 - Parent: connected to a node closer to the root or
 - Child: port connected to a node further from the root
- Any unconnected ports are labeled "off" and do not participate in further arbitration processes.
- Any loop in the topology is detected by a timeout in this process.

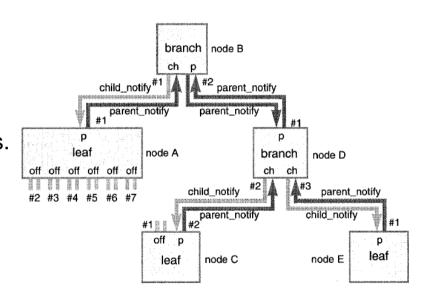


- The first step is for all leaf nodes to notify their probable parents. This is done by sending a parent_notify packet (signal)
- In this example, nodes A, C, and E send a parent_notify to their single connected port.
- This is the start of the parent-child handshake process.



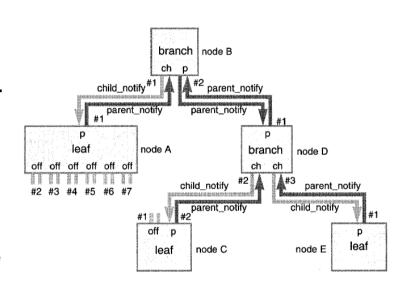


- The nodes internally recognize the parent_notify signals and mark the ports receiving them as child ports.
- If these nodes have one unidentified port, they send parent_notify up to their probable parents.
- At the same time the nodes send down child_notify signals to their child ports.
- Example here: Nodes B and D have one port remaining that is connected but not yet identified as child or parent. Therefore a parent_notify signal is sent.





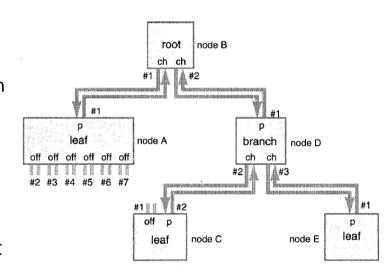
- When the leaf nodes receive the child_notify signal, that port is assigned as their parent port. Once assigned, their parent_notify signal is withdrawn.
- Nodes having only child ports are assigned as the root.
- It can occur that multiple nodes are receiving only parent_notify signals. This leads to a process called "root contention"
- Example:
 - Both nodes B and D discover that they are receiving the parent_notify port
 - Since one of the two nodes has to become the parent of the other, this collision of intentions starts a process called "root contention".
- This is resolved by withdrawing the parent_notify signals. A timer is assigned on each node with a random duration and the signals are resent after the time elapses.



Start of root contention



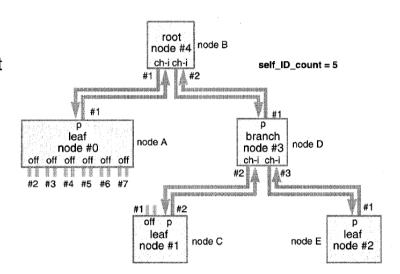
- When a node has all ports labeled as children, it takes the root function for itself.
- In this example:
 - Node D resends the signals first due to root contention
 - Node B has all ports as children and is assigned as the root.
- Note that the selection of the root node is not topology dependent. It is completely acceptable that the root node also be a leaf.
- The node that waits the longest after the bus reset to start participating in the tree identify process becomes the root.
- A particular node can be forced to wait a longer time by setting a force_root bit



3. Self-Identify Process

- The next step is to give each node a unique physical_ID
- The self-identify process uses a deterministic selection process, where the root node passes control of the media to the node attached to its lowest numbered connected port first. The child nodes then passes control in a recursive manner
- The first leaf node to receive control, during its self-identify process selects 0. The next node which receives control selects 1 and so on
- The root then passes control to the next lowest numbered port. When the nodes attached to all the ports of the root are finished, the root itself does a self-identify process.
- Note that each port of the node is individually numbered. There is no particular order to the numbering, it is just a way to give each port a unique label



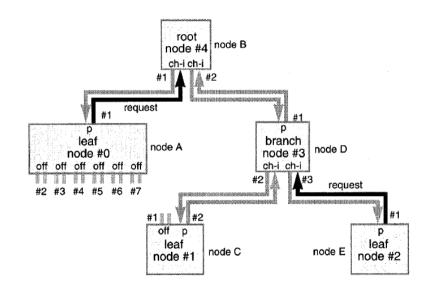


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Arbitration



- Once the self-identify process is complete, nodes can use the normal arbitration method to send packets
- Example:
 - Node A and E begin arbitrating at the same time by sending request to their parents

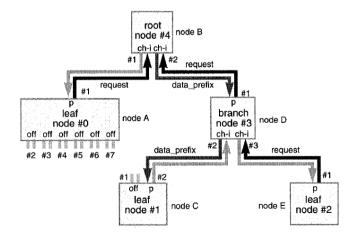


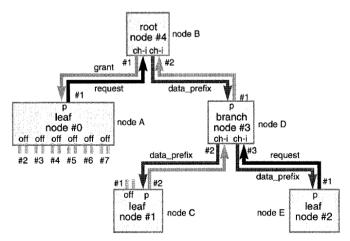
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Normal Arbitration



- Node D forwards the request and denies access to its other children by sending a *data prefix*.
- At the same time, node B receives request from node A and denies access to its other children.
- Since node B is root, it does not forward the request further and grants access to Node A.



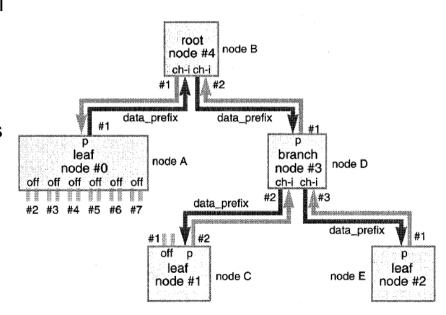


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Normal Arbitration



- The granted node sends a data_prefix signal to all other nodes to warn other nodes that data is about to be sent.
- When the parent (in the example it is the root) receives this data_prefix signal, it stops sending the grant.
- At this point, the physical connections between all the nodes are now in the same state and pointed away from the node that won the arbitration.



Fair Arbitration



- The normal cable and backplane arbitration methods guarantee that only one node will be transmitting at the end of the arbitration period.
- These methods only provide a strict priority access; the node with the highest natural priority will always win.
 - Example: For cable environment, closest node to the root always wins

Task 1: FireWire



10 min Time Allocated

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